AMENDMENTS TO THE CLAIMS

1   1.   (Previously Amended) A method to process commands in a computer

2   memory subsystem, comprising:

3   (a)   receiving a plurality of commands on a bus network connected to

4   said memory subsystem;

5   (b)   categorizing said received commands into command types;

6   (c)   placing each received command into a queue pertaining to its

7   respective command type;

8   (d)   determining memory cycle performance penalties of said

9   categorized commands in each of said queues;

10   (e)   reordering said categorized commands in each of said queues so

11   that one categorized command in each of said queues having the

12   least memory cycle performance penalty is selected for execution;

13   (f)   determining if each of said selected command is valid;

14   (g)   arbitrating said valid commands and selecting one of said valid

15   commands to execute; and

16   (h)   executing sequential valid commands of the same command type.

1   2.   (Original) The method of claim 1, wherein said command types are forms

2   of store and fetch operations.

1   3.   (Original) The method of claim 1, wherein said command types are

2   associated with a particular source or destination of said received

3   memory commands.

1   4.   (Original) The method of claim 3, wherein said particular source or

2   destination is a particular computer processor connected on said bus

3   network.

1   5.   (Original) The method of claim 3, wherein said particular source or

2   destination is a I/0 hub controller functionally connected on said bus

3   network.

1   6.   (Original) The method of claim 3, wherein said particular source or

2   destination is a switching fabric connected to said bus network.

1   7.   (Original) The method of claim 3, wherein said particular source or

2   destination is a compression/decompression engine functionally

3   connected to said bus network.

1   8.   (Original) The method of claim 1, wherein said command types which

2   originate from or are required for a particular application have priority.

1   9.   (Original) The method of claim 1, wherein said step of receiving a

2   plurality of commands further comprises determining if any of said

3   received commands have an address dependency and passing said

4   address dependency determination with said memory command.

1   10. (Original) The method of claim 1, wherein said step of determining

2   memory cycle performance penalties of said categorized commands

3   further comprises comparing a number of oldest received categorized

4   commands with each other.

1   11. (Original) The method of claim 9, wherein said step of determining

2   memory cycle performance penalties of said categorized commands

3   further comprises comparing a number of the oldest received categorized

4   commands with a currently chosen command.

1   12. (Original) The method of claim 9, wherein said step of determining

2   memory cycle performance penalties of said categorized commands

3   further comprises comparing a number of the oldest received categorized

4   commands with a previously chosen command.

1   13.   (Original) The method of claim 1, wherein said step of reordering said

2   categorized commands further comprises selecting the oldest of said

3   categorized commands that have the least memory cycle performance

4   penalty for execution.

1   14.   (Original) The method of claim 1, wherein said step of arbitrating said

2   reordered valid commands further comprises granting priority to said

3   type of command having said least memory cycle performance penalty.

1   15.   (Original) The method of claim 1, wherein said step of arbitrating said

2   reordered valid commands further comprises granting priority to a

3   command type other than said command type of said reordered valid

4   commands.

1   16.   (Previously Amended) The method of claim 1, wherein said step of

2   executing sequential valid commands of the same command type further

3   continues until a valid memory command of said command type is no

4   longer available, or until a predetermined number has been executed, or

5   until a memory command of another of said command types has higher

6   priority.

1   17.   (Previously Amended) A method to process commands in a computer

2   memory subsystem, comprising:

3   (a)   receiving a plurality of memory commands on a bus connected to

    said computer memory subsystem and determining the physical

5   location of the memory command in memory, and further

6   determining if any of said received memory commands have an

7   address dependency and passing said physical location and said

8   address dependency, if any, corresponding to said memory

9   command along with said memory command;

10   (b)   categorizing said received commands into command types based

11   on one of the following: STORE, FETCH, INTERVENTION STORE;

12   the source or destination of said received memory commands; the

13   program or application from which said memory commands

14   originate or are otherwise required;

15   (c)   determining memory cycle performance penalties of said

16   categorized commands by comparing a number of oldest received

17   categorized commands with each other, with a currently chosen

18   command, and with a previously chosen command;

19   (d)   reordering said categorized commands so that said categorized

20   commands having the least memory cycle performance penalty are

21   selected for execution and if more than one categorized command

22   has the least memory cycle performance penalty, then selecting the

23   oldest of said reordered commands for execution;

24   (e)   determining if said reordered commands are valid;

25   (f)   granting priority to said type of command having said least

26   memory   cycle performance penalty;

27   (g)   executing sequential valid commands of the same command type

28   until a valid command of the same type is not received or until a

29   predetermined number has been executed, or until a memory

30   command of another type has higher priority;

31   (h)   avoiding deadlock when an address dependency exists between

32   commands of different types by executing commands having the

33   command type of the oldest memory command.

1   18.   (Original) A method of processing memory commands in a computer

2   processing system having at least one command source on a bus

3   connected to a memory controller, said method comprising selecting a

4   memory command having the least memory cycle performance penalties

5   to execute and then executing a programmable number of other memory

6   commands of that type.

1   19.   (Currently Amended) A computer processing system, comprising:

2   (a)   a plurality of bus units, said bus units comprising at least one

3   computer processor, at least one I/O device; at least one memory

4   cache system connected to said at least one computer processor,

5   and at least one network communication device, said plurality of

6   bus units interconnected on a bus network, and said plurality of

7   bus units to issue memory commands, said memory commands

8   categorized into types;

9   (b)   at least one memory subsystem connected on a first bus to said

10   plurality of bus units, said memory subsystem responsive to said

11   memory commands and further comprising:

12   (i)   a memory controller connected to a command interface

13.   functionally connected to said first bus;

14   (ii)   a plurality of memory chips configured into memory banks;

15   said memory chips architected into memory cards attached

16   to at least one memory bus;

17   (iii)   a plurality of command FIFO queues, each of said command

18   FIFO queues associated with one of said command types into

19   which said memory commands are categorized;

20   (iv)   a plurality of comparison logic circuits, each of said plurality

21   of comparison logic circuits associated with each of said

22 plurality of command FIFO queues to determine which

23 memory commands of each of said command types have the

24 least memory cycle performance penalty <u>by comparing a</u>

25 <u>number of oldest received categorized commands with each</u>

26 <u>other, with a currently chosen command, and with a</u>

27 <u>previously chosen command;</u>

28 (v) an arbitration logic circuit to output said memory commands

29 of said determined command type having said least memory

30 cycle performance penalty to said plurality of memory chips.

1 20. (Original) The computer processing system of claim 19, wherein said

2 comparison logic circuit further determines the oldest of said memory

3. commands in each of said plurality of command FIFO queues.

1 21. (Currently Amended) A computer memory controller, comprising:

2 (a) means to receive a plurality of types of memory commands from a

3 plurality of command sources;

4 (b) means to determine the memory cycle performance penalty

5 associated with each memory command of each of said plurality of

6 types;

7   (c)   means to compare said memory commands of ~~one~~ <u>each</u> of said

8   types with other memory commands of the same type to determine

9   which of said memory commands have the least memory cycle

10   performance penalty;

11   (d)   means to compare said memory commands of ~~one~~ <u>each</u> of said

12   types with a current chosen memory command of the same type to

13   determine which of said memory commands have the least memory

14   cycle performance penalty;

15   (e)   means to compare said memory commands of ~~one~~ <u>each</u> of said

16   types with a previously chosen memory command of the same type

17   determine which of said memory commands have the memory

18   cycle performance penalty;

19   (f)   means to select one of said memory commands having the least

20   memory cycle performance penalty by selecting the oldest; and

21   (g)   means to continue execution of memory commands of the same

22   type as said selected memory command.